

Program Analysis

Call Graphs (Part 3)


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Slides adapted from Eric Bodder²²

Overview

- Introduction
- Single & efficient: CHA, RTA
- Analyzing assignments: VTA, DTA 
- Call graphs and points-to analysis:
Spark

Variable Type Analysis (VTA)

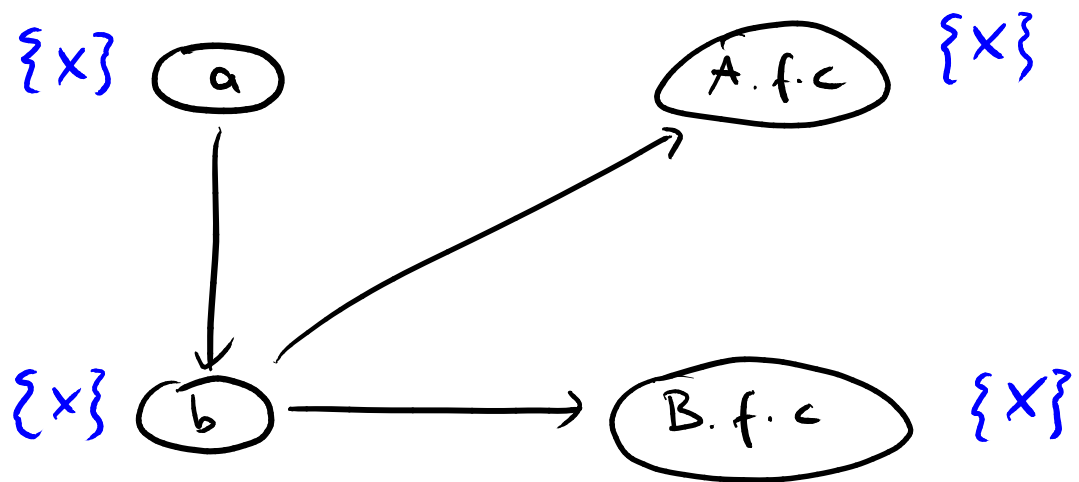
- Reason about **assignments**
- Infer what **types the objects involved in a call may have**
- Prune calls that are infeasible based on the inferred types

Example

```
a = new X();  
...  
b = a;  
...  
o.f(b);
```

```
public class A {  
    public void f(C c) {  
        c.m();  
    }  
}
```

```
public class B {  
    public void f(C c) {  
        c.m();  
    }  
}
```



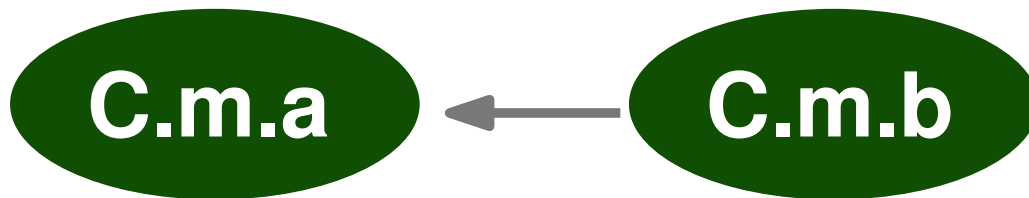
Type Propagation

Four steps:

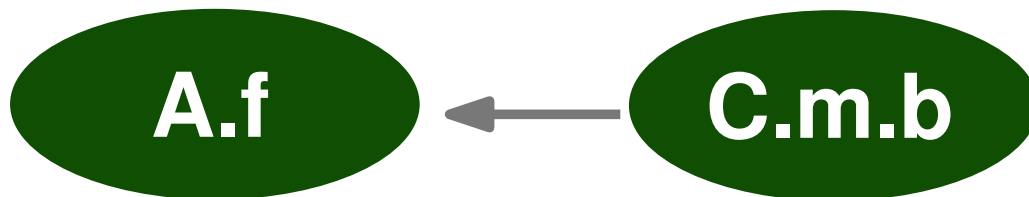
- Form **initial conservative call graph**
 - E.g., using CHA or RTA
- Build **type-propagation graph**
- **Collapse** strongly connected components
- **Propagate types** in one iteration

Building Type Propagation Graph

- Assume statement $a = b;$ is in method $C.m$



- Assume another statement $a.f = b;$ where field f is declared in A



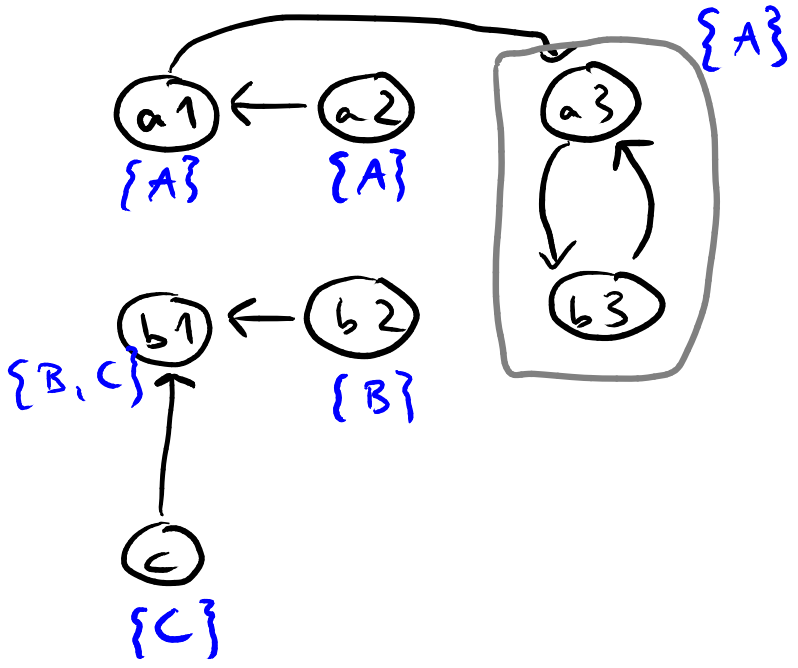
Example

```
A a1, a2, a3;  
B b1, b2, b3;  
C c;
```

```
a1 = new A();  
a2 = new A();  
b1 = new B();  
b2 = new B();  
c = new C();
```

```
a1 = a2;  
a3 = a1;  
a3 = b3;  
b3 = (B) a3;  
b1 = b2;  
b1 = c;
```


A
↑
B
↑
C



Note: Slide fixed w.r.t lecture video

Side Note: Field Representations

How does the analysis represent $a.f$?

- **Field-sensitive**: Represented as $a.f$
- **Field-insensitive**: Represented as $a.*$ or a
- **Field-based**: Represented as $A.f$, where A is class of a

Side Note: Field Representations

How does the analysis represent $a . f$?

- **Field-sensitive**: Represented as $a . f$
- **Field-insensitive**: Represented as $a . *$ or a
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VTA is field-based

Variable Type Analysis (VTA)

■ Pros

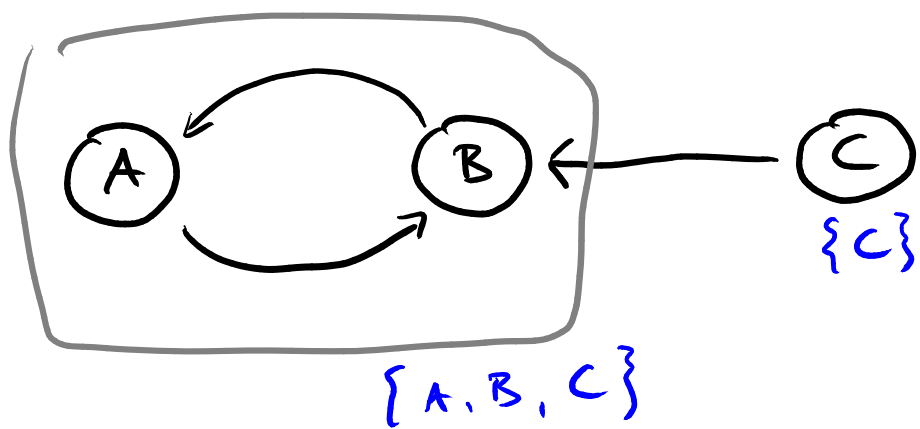
- **More precise than RTA**: Considers only those types that may actually reach the call site
- Still **relatively fast**

■ Cons

- **Requires initial call graph** (i.e., actually a refinement algorithm)
- **Some imprecision** remains, e.g., because of field-based analysis

Declared-Type Analysis (DTA)

- **“Small brother of VTA”**
- **Also reasons about assignments and how they propagate types**
- **But: Not per variable, but **per type****



Declared-Type Analysis (DTA)

■ Pros

- **Faster than VTA**: Graph is smaller, propagation is faster
- **More precise than RTA**

■ Cons

- **Less precise than VTA**: Does not distinguish variables of same type