

# **Program Analysis**

## **Call Graphs (Part 3)**

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Slides adapted from Eric Bodder<sup>22</sup>

# Overview

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- **Introduction**
- **Single & efficient: CHA, RTA**
- **Analyzing assignments: VTA, DTA** 
- **Call graphs and points-to analysis:  
Spark**

# Variable Type Analysis (VTA)

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- Reason about **assignments**
- Infer what **types the objects involved in a call may have**
- Prune calls that are infeasible based on the inferred types

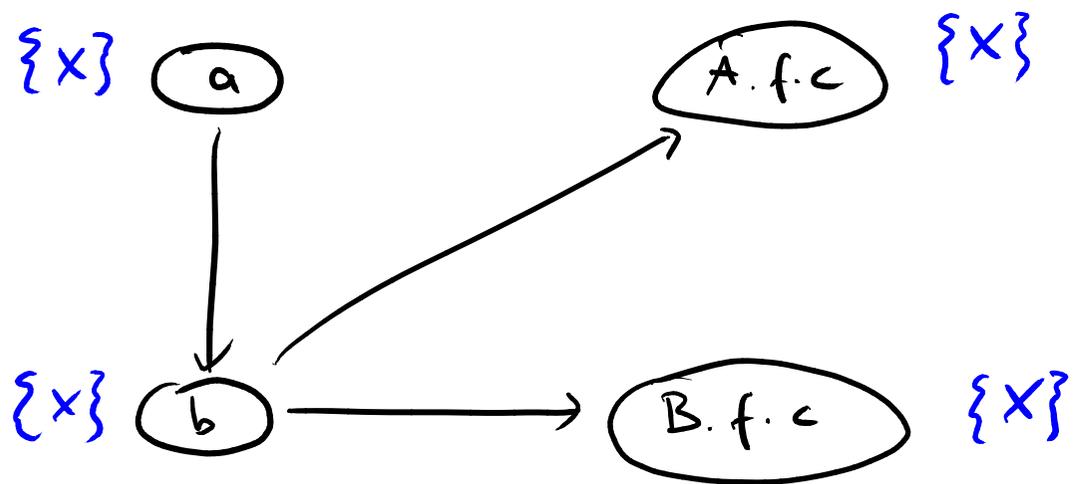
# Example

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```
a = new X();  
...  
b = a;  
...  
o.f(b);
```

```
public class A {  
    public void f(C c) {  
        c.m();  
    }  
}
```

```
public class B {  
    public void f(C c) {  
        c.m();  
    }  
}
```



# Type Propagation

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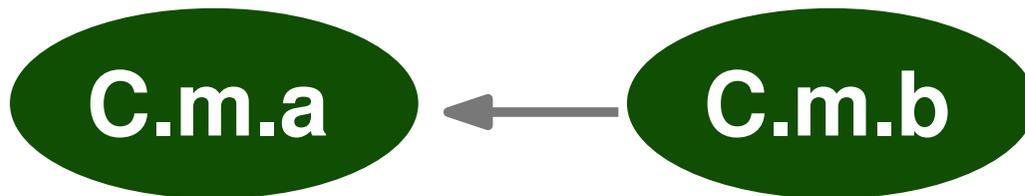
## Four steps:

- Form **initial conservative call graph**
  - E.g., using CHA or RTA
- Build **type-propagation graph**
- **Collapse** strongly connected components
- **Propagate types** in one iteration

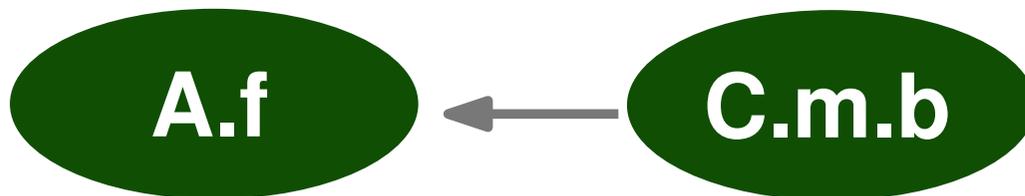
# Building Type Propagation Graph

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- Assume statement  $a = b;$  is in method  $C.m$



- Assume another statement  $a.f = b;$  where field  $f$  is declared in  $A$



# Example

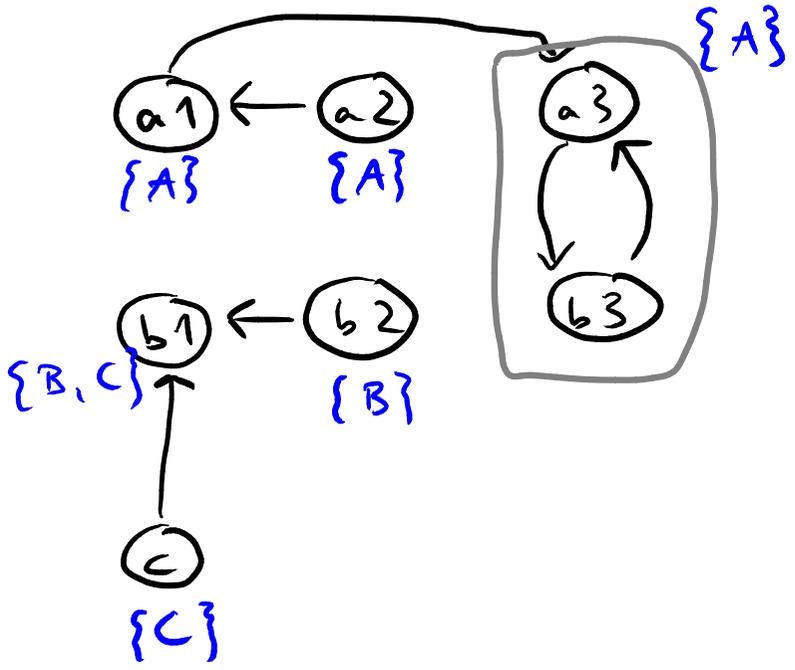
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```
A a1, a2, a3;  
B b1, b2, b3;  
C c;
```

```
a1 = new A();  
a2 = new A();  
b1 = new B();  
b2 = new B();  
c = new C();
```

```
a1 = a2;  
a3 = a1;  
a3 = b3;  
b3 = (B) a3;  
b1 = b2;  
b1 = c;
```

A  
↑  
B  
↑  
C



Note: Slide fixed w.r.t lecture video

# Side Note: Field Representations

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How does the analysis represent  $a.f$ ?

- **Field-sensitive**: Represented as  $a.f$
- **Field-insensitive**: Represented as  $a.*$  or  $a$
- **Field-based**: Represented as  $A.f$ , where  $A$  is class of  $a$

# Side Note: Field Representations

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**VTA is field-based**

# Variable Type Analysis (VTA)

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## ■ Pros

- **More precise than RTA**: Considers only those types that may actually reach the call site
- Still **relatively fast**

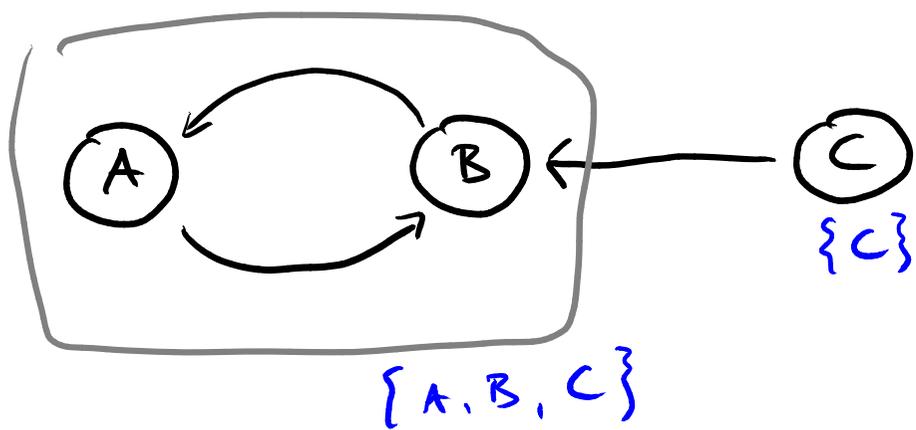
## ■ Cons

- **Requires initial call graph** (i.e., actually a refinement algorithm)
- **Some imprecision** remains, e.g., because of field-based analysis

# Declared-Type Analysis (DTA)

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- **“Small brother of VTA”**
- **Also reasons about assignments and how they propagate types**
- **But: Not per variable, but **per type****



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# Declared-Type Analysis (DTA)

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## ■ Pros

- **Faster than VTA**: Graph is smaller, propagation is faster
- **More precise than RTA**

## ■ Cons

- **Less precise than VTA**: Does not distinguish variables of same type