

Programming Paradigms

Lecture 4: Names, Scopes, and Bindings

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Names in PLs

Abstraction in two dimensions

- **From hardware**

- Variable names abstract away how exactly values are stored

- **From implemented functionality**

- Function names abstract from the implemented behavior

Binding

- Association between **entities** and their **names**, e.g.,
 - A variable bound to a memory object
 - A function bound to the code implementing the function
- Different languages have **different rules**
 - E.g., static vs. dynamic binding

Scope

- **Scope of a binding**: Textual region where binding is active
- **Scope**: Maximal region where no bindings change

Example (Python):

```
x = 1
def f():
    x = 2
    y = x
```

Scope

- **Scope of a binding**: Textual region where binding is active
- **Scope**: Maximal region where no bindings change

Example (Python):

```
x = 1          ] Outer scope
def f():
    x = 2      ] Scope of
    y = x      ] function
```

Overview

- **Object lifetime and storage management**
- **Scopes**
- **Aliasing and overloading**
- **Binding of referencing environments**

Object Lifetime

Every **memory object** has a **lifetime**

- Global variables: Entire program execution
- Local variables: Function execution

Object lifetime vs. binding lifetime

- A single object may be bound to multiple names
- Bindings may be concurrent

Example 1

→
program execution

→
value 3

→
x →
 y

```
fun f() {  
  x = 3  
  return x  
}
```

y = f()

Example 2:

→
program execution

→
object

→
binding

↳ usually a bug ("dangling reference")

↳ use-after-free attack in C

Storage Allocation

Three kinds of **memory objects**

- **Static**

- Absolute address retained throughout execution

- **Stack**

- Usually within subroutines
- Allocation/deallocation on call/return

- **Heap**

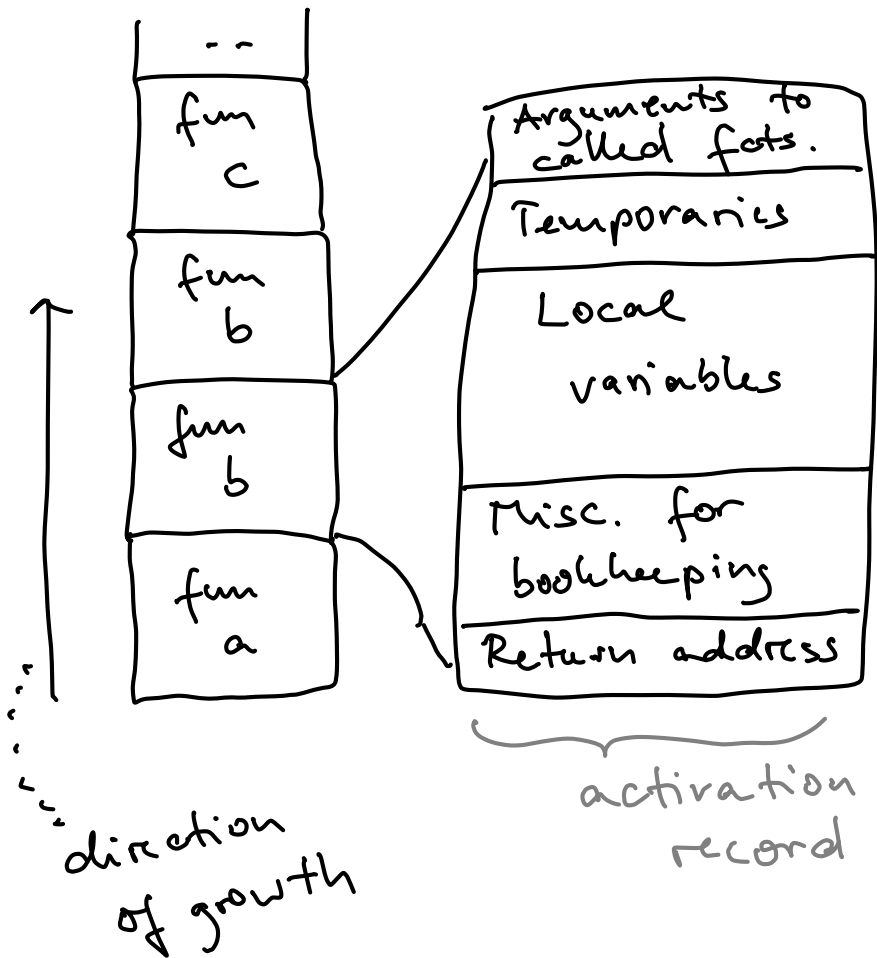
- Allocation and deallocation at arbitrary times

Statically Allocated Memory

Depending on the PL, used, e.g., for

- Global variables
- Constant literals
- Symbol tables
- Program code itself
- Compile-time constants
 - Even if local to function

Stack-based Allocation



```

fun c() {
    ...
}
fun b(...) {
    if ...
        b()
    else
        c()
}
fun a() {
    b()
}
// main
a()

```

Heap-based Allocation

- **For dynamically allocated data structures and objects whose size is statically unknown**
 - E.g., objects in Java
- **Some PLs: Managed memory**
 - Unreachable objects: Implicitly deallocated
 - Unreachable = No active binding
 - Less control but fewer bugs
 - E.g., no use-after-free

Quiz: Memory Allocation

Where are the following data objects stored (Java)?

- The integer 23
- The string "John"
- The Person object
- The reference variable p

```
class Person {
    int pid;
    String name;

    // constructor
}

public class Driver {
    public static void main(String[] args) {
        int id = 23;
        String pName = "John";
        Person p = null;
        p = new Person(id, pName);
    }
}
```

Quiz: Memory Allocation

```
class Person {  
    int pid;  
    String name;  
  
    // constructor  
}
```

```
public class Driver {  
    public static void main(String[] args) {  
        int id = 23;  
        String pName = "John";  
        Person p = null;  
        p = new Person(id, pName);  
    }  
}
```

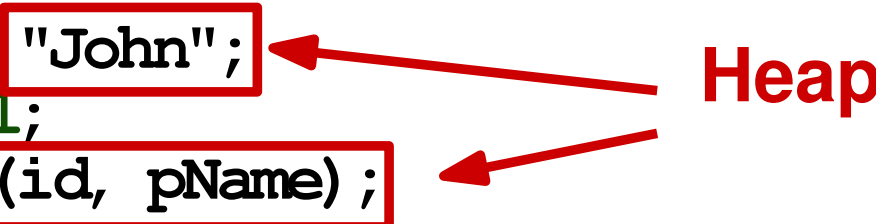
**Stack (in allocation
frame of main)**

Quiz: Memory Allocation

```
class Person {
    int pid;
    String name;

    // constructor
}

public class Driver {
    public static void main(String[] args) {
        int id = 23;
        String pName = "John";
        Person p = null;
        p = new Person(id, pName);
    }
}
```



The diagram illustrates memory allocation for the provided code. Two red boxes highlight the string "John" and the expression "new Person(id, pName);". Red arrows point from the word "Heap" to both of these boxes, indicating that both the string and the object are allocated on the heap.