

Programming Paradigms

Names, Scopes, and Bindings

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Names in PLs

Abstraction in two dimensions

- **From hardware**

- Variable names abstract away how exactly values are stored

- **From implemented functionality**

- Function names abstract from the implemented behavior

Binding

- Association between **entities** and their **names**, e.g.,
 - A variable bound to a memory object
 - A function bound to the code implementing the function
- Different languages have **different rules**
 - E.g., static vs. dynamic binding

Scope

- **Scope of a binding**: Textual region where binding is active
- **Scope**: Maximal region where no bindings change

Example (Python):

```
x = 1
def f():
    x = 2
    y = x
```

Scope

- **Scope of a binding**: Textual region where binding is active
- **Scope**: Maximal region where no bindings change

Example (Python):

```
x = 1           ] Outer scope
def f():
    x = 2       ] Scope of
                ] function
    y = x
```

Overview

- **Object lifetime and storage management**
- **Scopes**
- **Aliasing and overloading**
- **Binding of referencing environments**

Object Lifetime

Every **memory object** has a **lifetime**

- Global variables: Entire program execution
- Local variables: Function execution

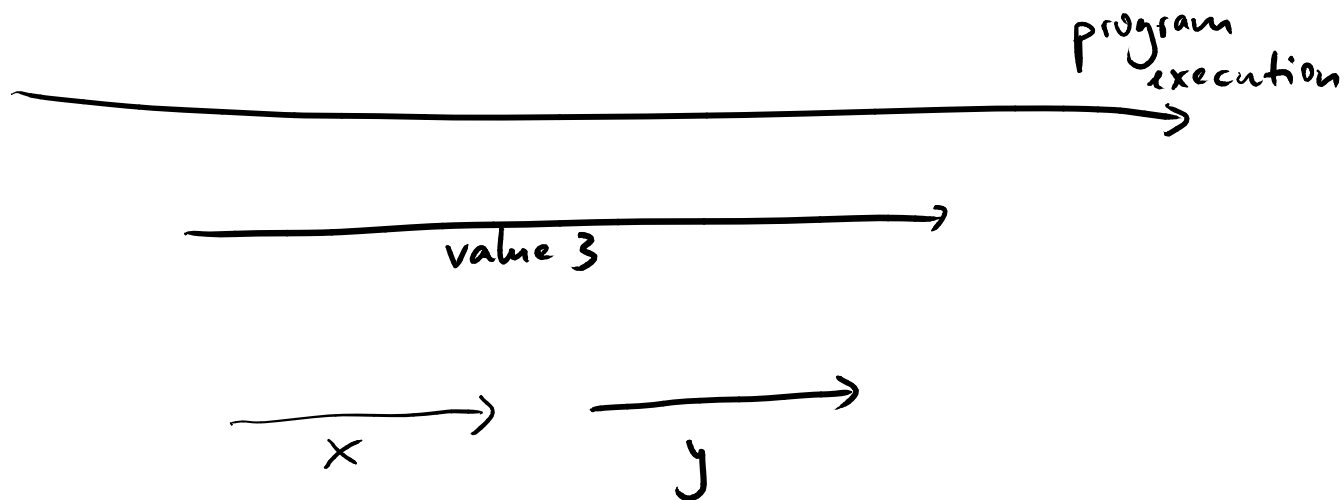
Object lifetime vs. binding lifetime

- A single object may be bound to multiple names
- Bindings may be concurrent

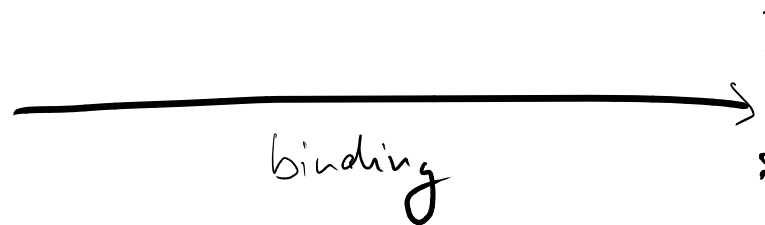
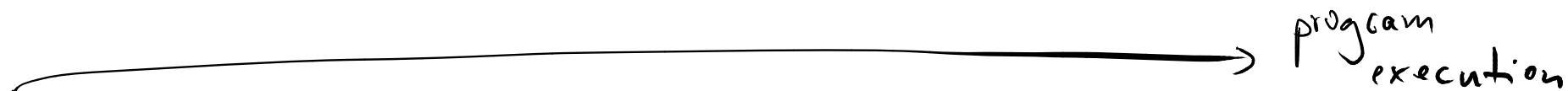
Example 1

```
fun f() {  
  x = 3  
  return x  
}
```

$y = f()$



Example 2



usually a bug ("dangling reference")
↳ use-after-free attack in C

Storage Allocation

Three kinds of **memory objects**

- **Static**

- Absolute address retained throughout execution

- **Stack**

- Usually within subroutines
- Allocation/deallocation on call/return

- **Heap**

- Allocation and deallocation at arbitrary times

Statically Allocated Memory

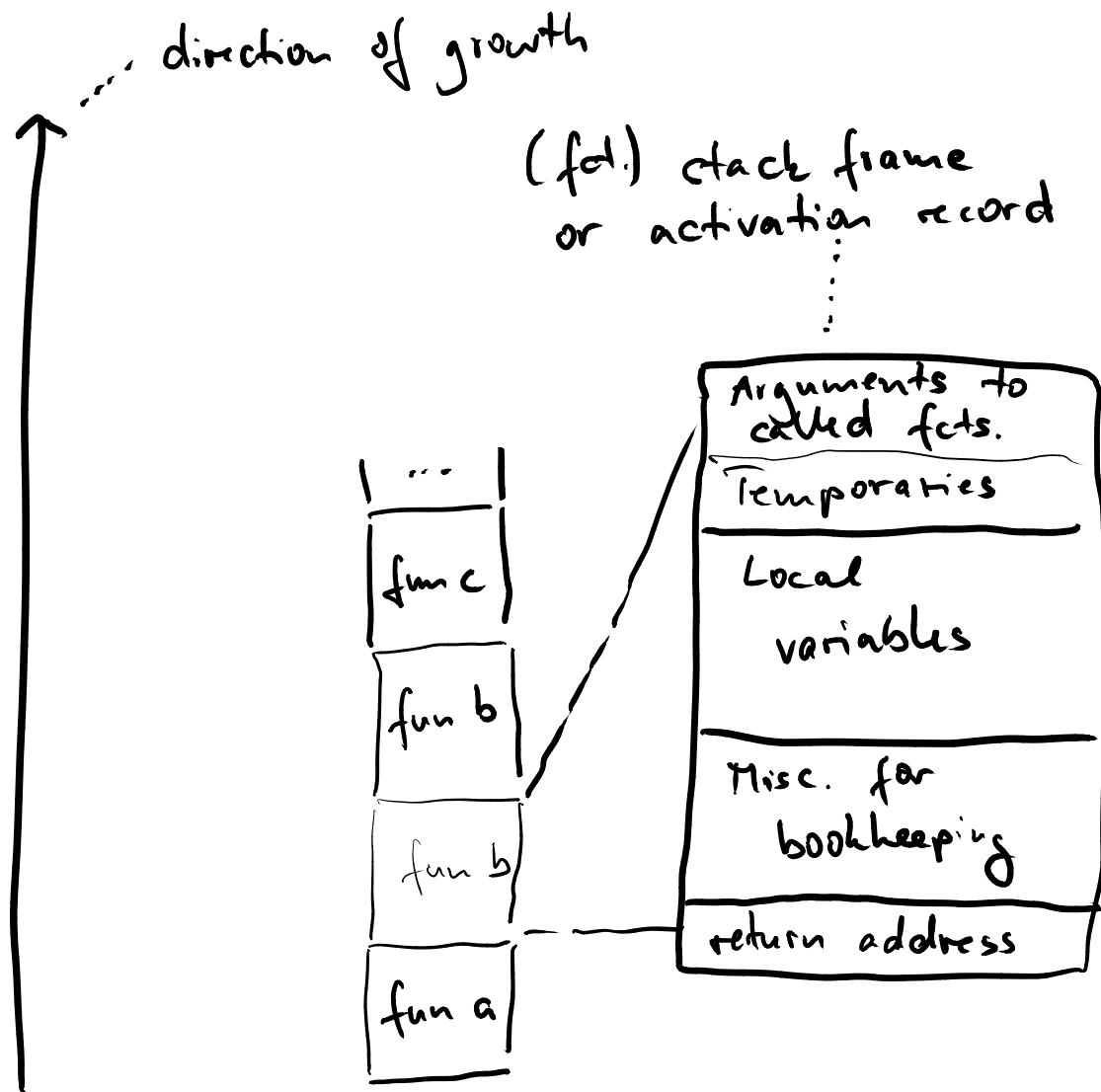
Depending on the PL, used, e.g., for

- Global variables
- Constant literals
- Symbol tables
- Program code itself
- Compile-time constants
 - Even if local to function

Stack-based Allocation

```

fun c() {
    ...
}
fun b() {
    if ...
        b()
    else
        c()
}
fun a() {
    b()
}
// main
a()
  
```



Heap-based Allocation

- **For dynamically allocated data structures and objects whose size is statically unknown**
 - E.g., objects in Java
- **Some PLs: Managed memory**
 - Unreachable objects: Implicitly deallocated
 - Unreachable = No active binding
 - Less control but fewer bugs
 - E.g., no use-after-free

Quiz: Memory Allocation

Where are the following data objects stored (in Java)?

- The integer 4
- The reference variable `c`
- The `Course` object

```
class Course {  
    String name;  
    int credits;  
  
    // constructor  
}
```

```
public class App {  
    public static void main(String[] args) {  
        String name = "PP";  
        int credits = 3+1;  
  
        Course c = null;  
        c = new Course(name, credits);  
    }  
}
```

Quiz: Memory Allocation

```
class Course {  
    String name;  
    int credits;  
  
    // constructor  
}
```

```
public class App {  
    public static void main(String[] args) {  
        String name = "PP";  
        int credits = 3+1; ← Stack (in allocation  
  
        Course c = null; ← frame of main)  
        c = new Course(name, credits);  
    }  
}
```

Quiz: Memory Allocation

```
class Course {
    String name;
    int credits;

    // constructor
}

public class App {
    public static void main(String[] args) {
        String name = "PP";
        int credits = 3+1;

        Course c = null;
        c = new Course(name, credits); ← Heap
    }
}
```


Quiz: Memory Allocation

```
class Course {  
    String name;  
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public class App {  
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        String name = "PP";  
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        c = new Course(name, credits);  
    }  
}
```

Bonus: Where is the string stored?

Quiz: Memory Allocation


```
class Course {  
    String name;  
    int credits;  
  
    // constructor  
}
```

```
public class App {  
    public static void main(String[] args) {  
        String name = "PP";  
        int credits = 3+1;  
  
        Course c = null;  
        c = new Course(name, credits);  
    }  
}
```

**Bonus: Where
is the string
stored?**

**String pool in
heap space**

Overview

- **Object lifetime and storage management**
- **Scopes** 
- **Aliasing and overloading**
- **Binding of referencing environments**

Scoping Rules

- **Scoping rules:** Define which bindings are active
 - I.e., what's the **meaning of a name** at a given **program point**?
- **Each PL defines its scoping rules**
 - E.g., Basic has only one scope
 - Most PLs have nested scopes for subroutines

Nested Scopes

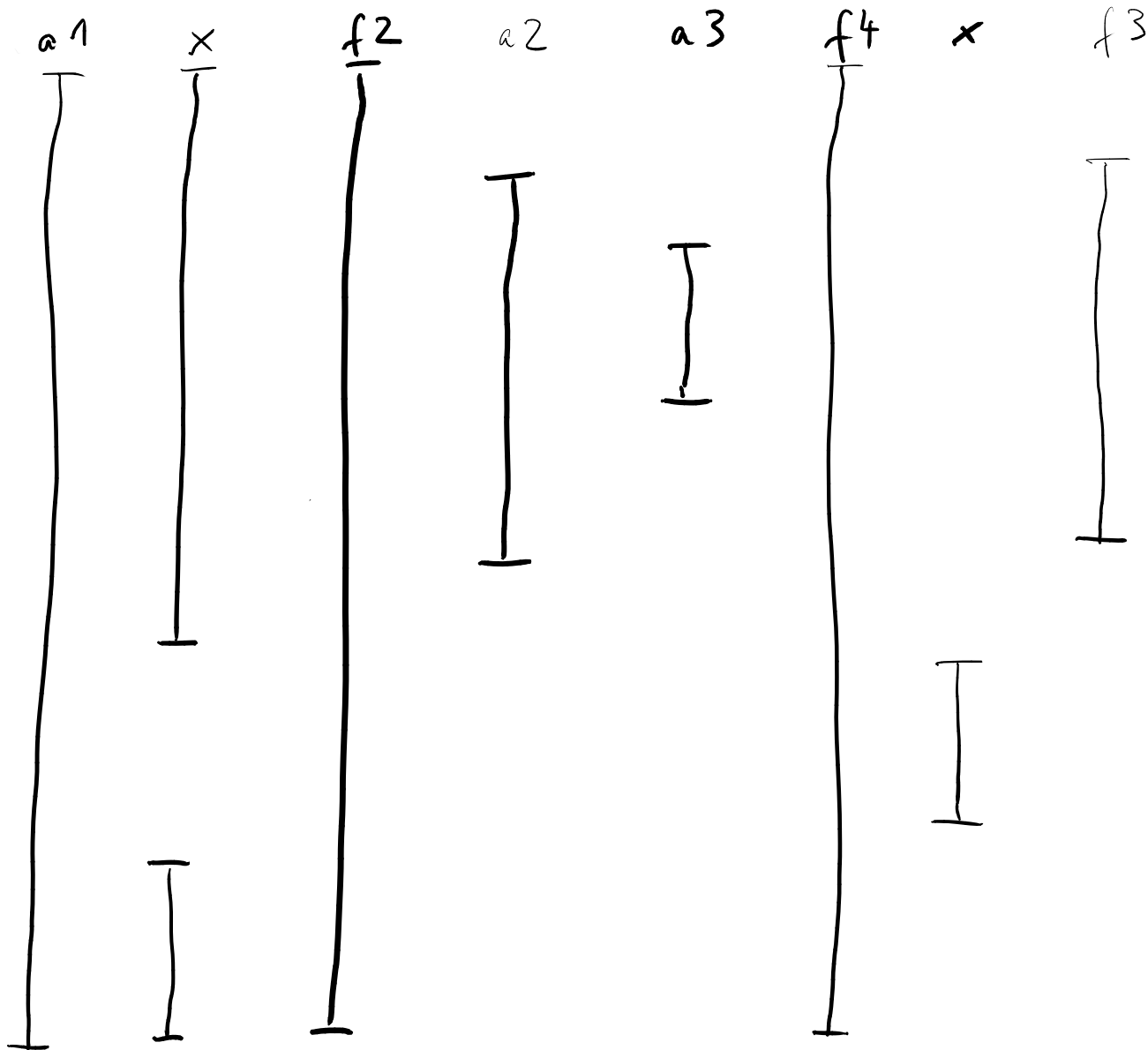
- Common for **nested subroutines**
- Each **subroutine** has its **own scope**
- **Closest nested scope rule**
 - Name is known in scope where it is declared and all scopes nested within
 - Inner scopes can hide names from outer scopes

Example

```

fun f1 (a1) {
  var x
  fun f2 (a2) {
    fun f3 (a3) {
      :
    }
    :
  }
  fun f4 (a4) {
    fun f5 (a5) {
      var x
    }
  }
}

```



Static vs. Dynamic Scoping

Static scoping

- Binding of a name can be derived from program text
- Most common in today's PLs

Dynamic scoping

- Binding of a name depends on control flow
 - I.e., not known statically (in general)

Example

Pseudo code:

```
global x = 1
fun a() {
    local x = 3
    b()
}
fun b() {
    y = x
}
a()
```


Example

Pseudo code:

```
global x = 1
fun a() {
    local x = 3
    b()
}
fun b() {
    y = x
}
a()
```

Static scoping:

y gets value 1 because

- b doesn't have a local variable called x
- the surrounding static scope provides the global variable x

Example

Pseudo code:

```
global x = 1
fun a() {
    local x = 3
    b()
}
fun b() {
    y = x
}
a()
```

Dynamic scoping:

y gets value 3 because

- b doesn't have a local variable called x
- the dynamically closest scope provides the local variable x of a

Quiz: Dynamic Scoping

What does this Perl code print?

(Hint: Perl uses dynamic scoping for `local` variables)

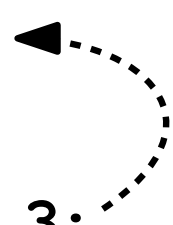
```
$b = 5;
sub foo {
    return $b;
}
sub bar {
    local $b = 3;
    return foo();
}
print bar();
```

Quiz: Dynamic Scoping

What does this Perl code print?

(Hint: Perl uses dynamic scoping for `local` variables)

```
$b = 5;
sub foo {
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sub bar {
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print bar();
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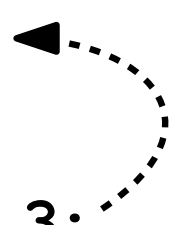
**Scope of `local $b`
dynamically extends
into invocation of `foo`**

Quiz: Dynamic Scoping

What does this Perl code print?

(Hint: Perl uses dynamic scoping for `local` variables)

```
$b = 5;
sub foo {
    return $b;
}
sub bar {
    local $b = 3;
    return foo();
}
print bar();
```



**Scope of `local $b`
dynamically extends
into invocation of `foo`**

Answer: 3

Quiz: Static Scoping

**What does this Python code print?
(Hint: Python uses static scoping)**


```
x = "l"
def f():
    def g():
        print(x)
    def h():
        g()
        x = "s"
        print(x)
    x = "e"
    h()
    print(x)
f()
print(x)
```

Quiz: Static Scoping

**What does this Python code print?
(Hint: Python uses static scoping)**

```
x = "l"
def f():
    def g():
        print(x) # (1) x in f : "e"
    def h():
        g()
        x = "s"
        print(x) # (2) x in h : "s"
    x = "e"
    h()
    print(x) # (3) x in f : "e"
f()
print(x) # (4) x in main: "l"
```

Function Stack vs. Static Scopes

- 
- Push allocation frames on calls
 - Pop frames on returns

- Not affected by which functions get called

How to **resolve bindings** outside of current scope?

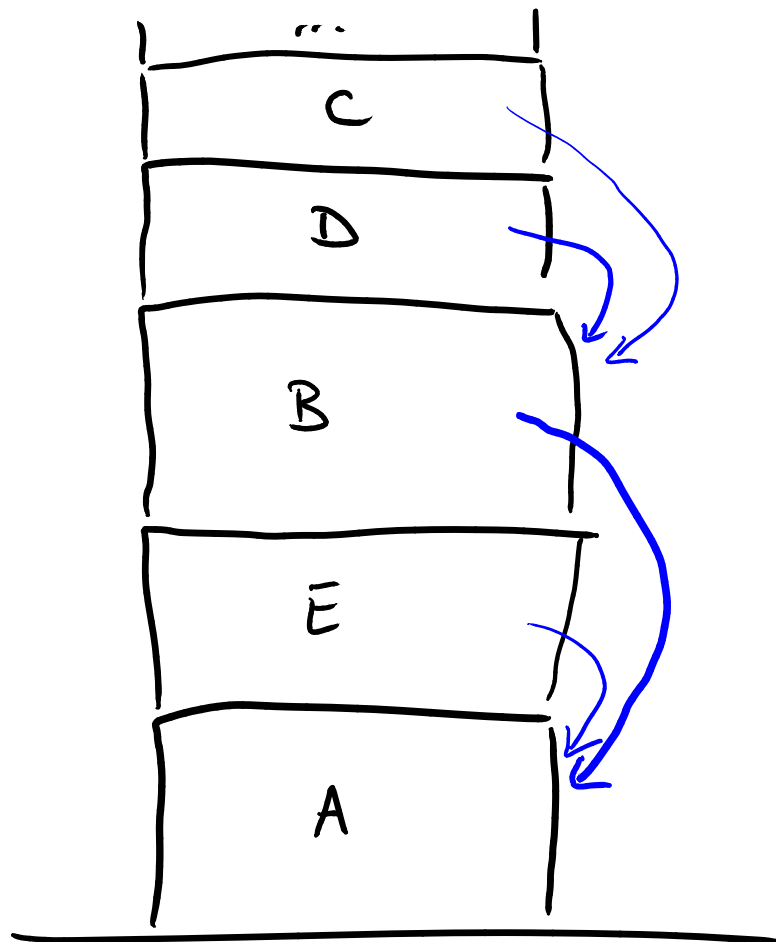
- Each allocation frame has a **static link** to its parent scope

Example

Nested functions



Function stack



→ ... static
link

Built-in Objects

Many PLs have **built-in (or predefined) objects**

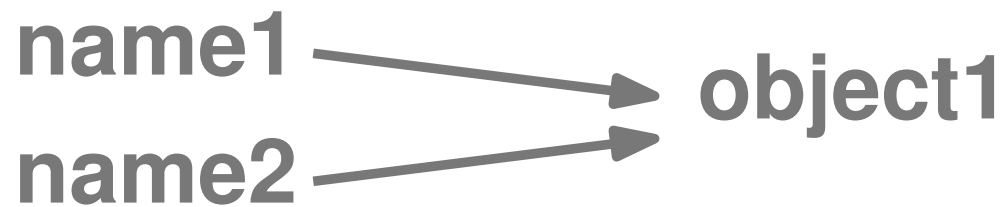
- E.g., for built-in types and APIs
- Invisible, **outer-most scope**
- Accessible from all scopes, except if hidden

Overview

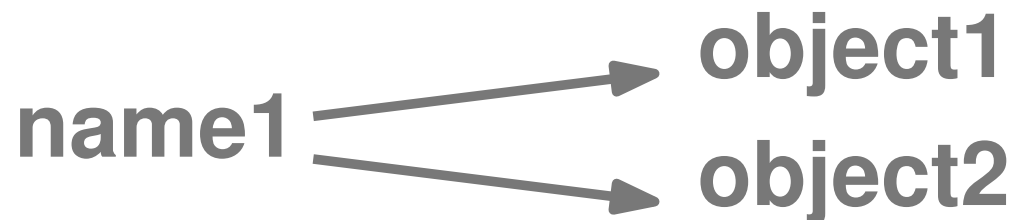
- **Object lifetime and storage management**
- **Scopes**
- **Aliasing and overloading** ←
- **Binding of referencing environments**

Aliasing and Overloading

Aliasing: Two more more names refer to the same object



Overloading: A name refers to two more objects



Aliasing: Example

```
#include <stdio.h>
```

```
void half(double& a)
{ // argument passed by reference
  a = a / 2;
}
```

```
int main(int argc, const char* argv[])
{
  double n = 5.0;
  double *p = &n; // pointer to value stored in n

  half(n);
  half(*p);

  printf("%f\n", n);
}
```

Aliasing: Example

```
#include <stdio.h>
```

```
void half(double& a)
{ // argument passed by reference
  a = a / 2;
}
```

```
int main(int argc, const char* argv[])
{
```

```
  double n = 5.0;
```

```
  double *p = &n; // pointer to value stored in n
```

```
  half(n);
```

```
  half(*p);
```

```
  printf("%f\n", n);
```

```
}
```

**Aliases to same
memory object**

Result: 1.250000

Overloading: Example

```
class Overloading{
    void foo() {}
    void foo(int n) {}
    void foo(String s) {}

    public static void main(String[] args) {
        Overloading o = new Overloading();
        o.foo(...);
    }
}
```

Overloading: Example

```
class Overloading{
```

```
void foo() {}  
void foo(int n) {}  
void foo(String s) {}
```

Three methods,
all with name
“foo”

```
public static void main(String[] args) {  
    Overloading o = new Overloading();  
    o.foo(...);
```



Resolution of name
depends on arguments

Overview

- **Object lifetime and storage management**
- **Scopes**
- **Aliasing and overloading**
- **Binding of referencing environments** ←

Referencing Environment

Complete set of **bindings at a point** in the program

- Determined by scoping rules (e.g., static or dynamic scoping)

What if we create a **reference to a function**?

- **When to apply** the scoping rules?

Example

Pseudo code:

```
function a() {  
  var x = 23;  
  function b() {  
    console.log(x);  
  }  
  x = 42;  
  return b;  
}  
b = a();  
var x = 5;  
b();
```

Example

Pseudo code:

```
function a() {  
  var x = 23;  
  function b() {  
    console.log(x);  
  }  
  x = 42;  
  return b;  
}  
b = a();  
var x = 5;  
b();
```

**Reference to
a function**



Example

Pseudo code:

```
function a() {  
  var x = 23;  
  function b() {  
    console.log(x);  
  }  
  x = 42;  
  return b;  
}  
b = a();  
var x = 5;  
b();
```

**Reference to
a function**

**Function
called here**

Example

Pseudo code:

```
function a() {  
  var x = 23;  
  function b() {  
    console.log(x);  
  }  
  x = 42;  
  return b;  
}  
b = a();  
var x = 5;  
b();
```

**What memory object
is **x** bound to?**

Shallow Binding

Referencing environment created **when function is called**

- Common in languages with **dynamic scoping**

Shallow Binding

Referencing environment created **when function is called**

- Common in languages with **dynamic scoping**

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function a() {  
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    console.log(x);  
  }  
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b = a();  
var x = 5;  
b();
```


Shallow Binding

Referencing environment created **when function is called**

- Common in languages with **dynamic scoping**

```
function a() {  
  var x = 23;  
  function b() {  
    console.log(x);  
  }  
  x = 42;  
  return b;  
}  
b = a();  
var x = 5;  
b();
```

**x bound to the global
variable initialized to 5;
code prints 5**

Deep Binding

Referencing environment created **when the reference to the function is created**

- Common in languages with **static scoping**

Deep Binding

Referencing environment created **when the reference to the function is created**

- Common in languages with **static scoping**

```
function a() {  
  var x = 23;  
  function b() {  
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  }  
  x = 42;  
  return b;  
}  
b = a();  
var x = 5;  
b();
```

Deep Binding

Referencing environment created **when the reference to the function is created**

- Common in languages with **static scoping**

```
function a() {  
  var x = 23;  
  function b() {  
    console.log(x);  
  }  
  x = 42; ←  
  return b; ←  
}  
b = a();  
var x = 5;  
b();
```

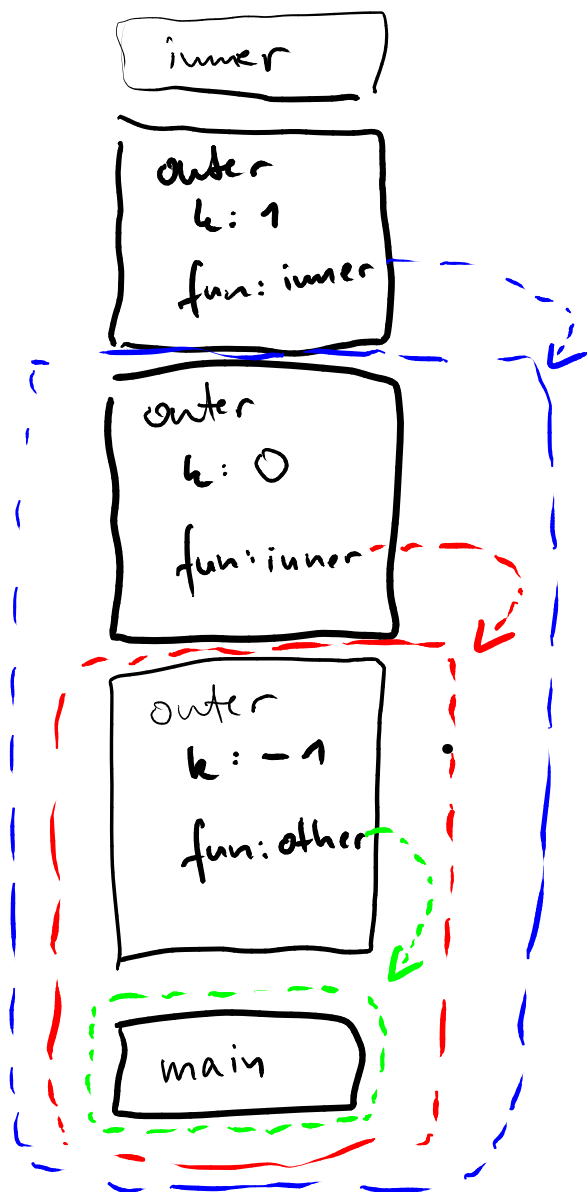
x bound to the local variable initialized to 23; code prints 42, as this is the most recent value of x

Closure

- Implementation of deep binding
- Closure = Representation of referencing environment + function itself
- When creating reference to function, closure is created

Example: Closures

```
function outer(k, fun) {  
    function inner() {  
        console.log(k);  
    }  
  
    if (k > 0)  
        fun();  
    else  
        outer(k + 1, inner)  
}  
  
function other() {}  
  
outer(-1, other);
```



-----> } referencing
 -----> } environments
 -----> } captured by
 } closures

prints 0

Quiz: Scopes and Bindings

Which of the following statements is true?

- Heap-allocated memory is freed at the end of each function call.
- A memory object is bound to at most one name.
- Built-in objects are available in an invisible, outer-most scope.
- With shallow binding, the referencing environment is created when a function is called.

Quiz: Scopes and Bindings

Which of the following statements is true?

- ~~Heap-allocated memory is freed at the end of each function call.~~
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- Built-in objects are available in an invisible, outer-most scope.
- With shallow binding, the referencing environment is created when a function is called.

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