Programming Paradigms Names, Scopes, and Bindings

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Software Lab, University of Stuttgart Summer 2022

Names in PLs

Abstraction in two dimensions

From hardware

 Variable names abstract away how exactly values are stored

From implemented functionality

 Function names abstract from the implemented behavior

Binding

Association between entities and their names, e.g.,

- □ A variable bound to a memory object
- A function bound to the code implementing the function
- Different languages have different rules

□ E.g., static vs. dynamic binding



- Scope of a binding: Textual region where binding is active
- Scope: Maximal region where no bindings change

```
Example (Python):
```

```
x = 1
def f():
x = 2
y = x
```



- Scope of a binding: Textual region where binding is active
- Scope: Maximal region where no bindings change

Example (Python):

- x = 1 **Outer scope**
- def f(): x = 2 **Scope of**
 - y = x J function

Overview

- Object lifetime and storage management
- Scopes
- Aliasing and overloading
- Binding of referencing environments

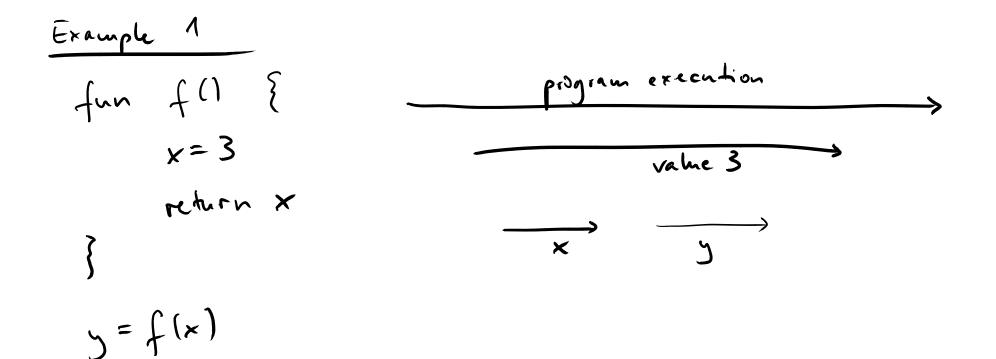
Object Lifetime

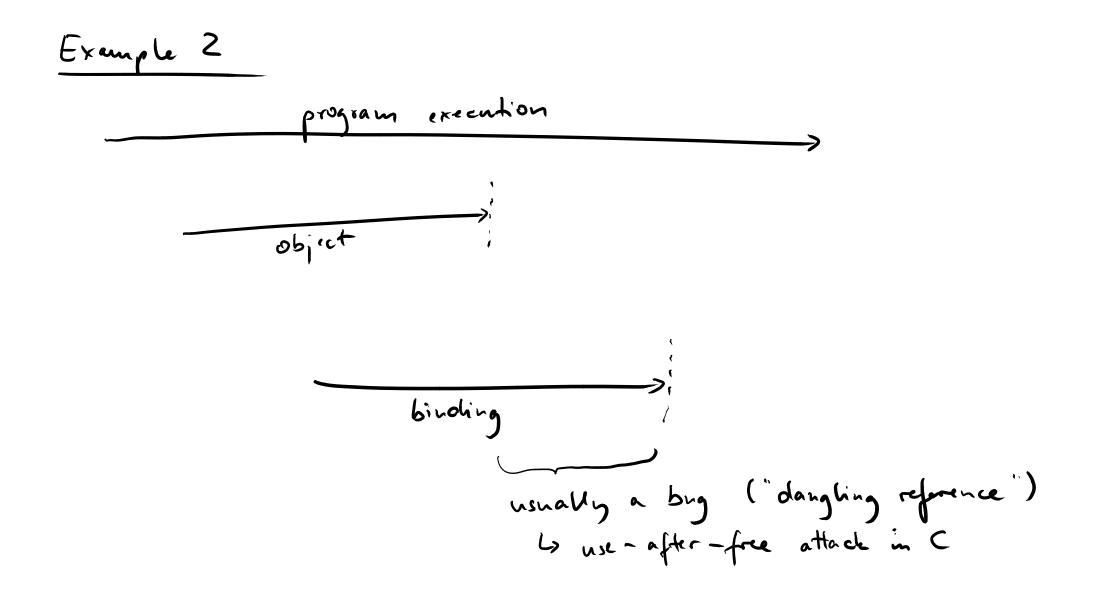
Every memory object has a lifetime

- Global variables: Entire program execution
- Local variables: Function execution

Object lifetime vs. binding lifetime

- A single object may be bound to multiple names
- Bindings may be concurrent





Storage Allocation

Three kinds of memory objects

- Static
 - □ Absolute address retained throughout execution
- Stack
 - Usually within subroutines
 - Allocation/deallocation on call/return

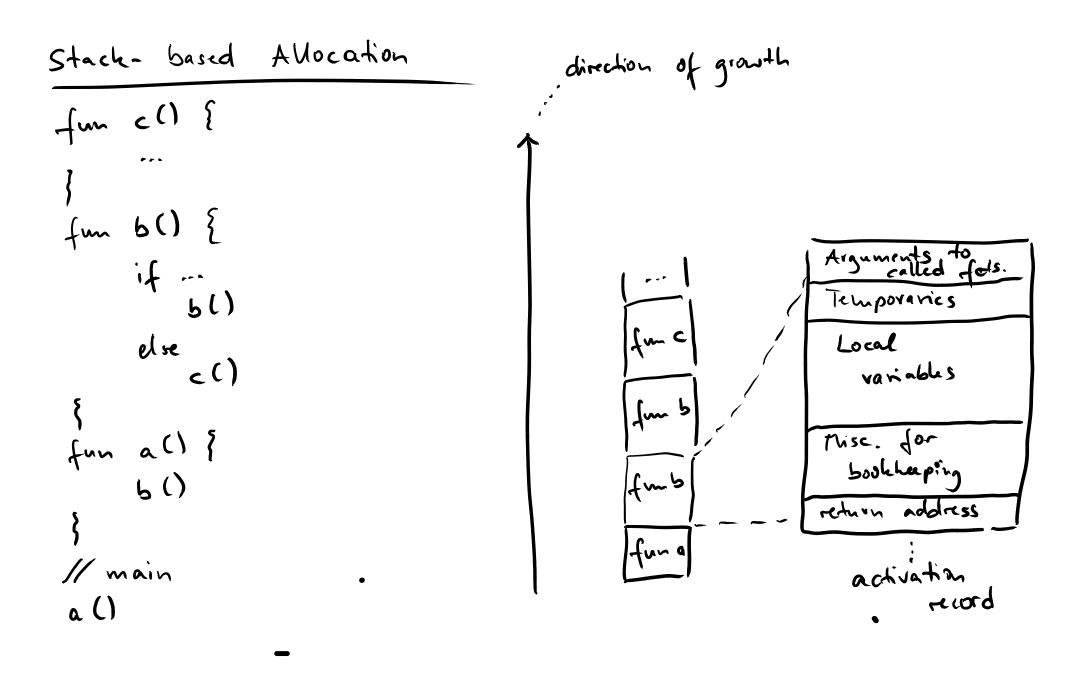
Heap

Allocation and deallocation at arbitrary times

Statically Allocated Memory

Depending on the PL, used, e.g., for

- Global variables
- Constant literals
- Symbol tables
- Program code itself
- Compile-time constants
 - □ Even if local to function



Heap-based Allocation

For dynamically allocated data structures and objects whose size is statically unknown

E.g., objects in Java

Some PLs: Managed memory

- Unreachable objects: Implicitly deallocated
 - Unreachable = No active binding
- □ Less control but fewer bugs
 - E.g., no use-after-free

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```
class Course {
   String name;
   int credits;
```

```
// constructor
}
```

Where are the following data objects stored (in Java)?

- The integer 4
- **The reference variable** c
- The Course object

```
public class App {
    public static void main(String[] args) {
        String name = "PP";
        int credits = 3+1;
        Course c = null;
        c = new Course(name, credits);
    }
```

```
class Course {
    String name;
    int credits;
    // constructor
}
public class App {
    public static void main(String[] args) {
        String name = "PP";
        int credits = 3+1;
                                      Stack (in allocation
        Course c = null;
                                    _ frame of main)
        c = new Course (name, credits);
```

```
class Course {
    String name;
    int credits;
    // constructor
}
public class App {
    public static void main(String[] args) {
        String name = "PP";
        int credits = 3+1;
        Course c = null;
                                             Heap
        c = new Course (name, credits);
```

```
class Course {
    String name;
    int credits;
    // constructor
}
public class App {
    public static void main(String[] args) {
        String name = "PP";
                                        Bonus: Where
        int credits = 3+1;
                                        is the string
        Course c = null;
        c = new Course (name, credits); stored?
    }
```

```
class Course {
    String name;
    int credits;
    // constructor
}
public class App {
    public static void main(String[] args) {
        String name = "PP";
                                        Bonus: Where
        int credits = 3+1;
                                        is the string
        Course c = null;
        c = new Course (name, credits); stored?
                                       String pool in
                                        heap space
                                                        14 - 5
```

Overview

- Object lifetime and storage management
- Scopes
- Aliasing and overloading
- Binding of referencing environments

Scoping Rules

- Scoping rules: Define which bindings are active
 - I.e., what's the meaning of a name at a given program point?

Each PL defines its scoping rules

□ E.g., Basic has only one scope

Most PLs have nested scopes for subroutines

Nested Scopes

- Common for nested subroutines
- Each subroutine has its own scope
- Closest nested scope rule
 - Name is known in scope where it is declared and all scopes nested within
 - Inner scopes can hide names from outer scopes

•

Static vs. Dynamic Scoping

Static scoping

- Binding of a name can be derived from program text
- Most common in today's PLs

Dynamic scoping

Binding of a name depends on control flow

□ I.e., not known statically (in general)

Example

Pseudo code:

```
global x = 1
fun a() {
    local x = 3
    b()
}
fun b() {
    y = x
}
a()
```

Example

Pseudo code:

```
global x = 1
fun a() {
    local x = 3
    b()
}
fun b() {
    y = x
}
a()
```

Static scoping:

y gets value 1 because b doesn't have a local variable called x and the surrounding static scope provides the global variable x

Example

Pseudo code:

```
global x = 1
fun a() {
    local x = 3
    b()
}
fun b() {
    y = x
}
a()
```

Dynamic scoping: y gets value 3 because b doesn't have a local variable called x and the dynamically closest scope provides the local variable x of a

Quiz: Dynamic Scoping

What does this Perl code print? (Hint: Perl uses dynamic scoping for local variables)

```
$b = 0;
sub foo {
    return $b;
}
sub bar {
    local $b = 1;
    return foo();
}
print bar();
```

Quiz: Dynamic Scoping

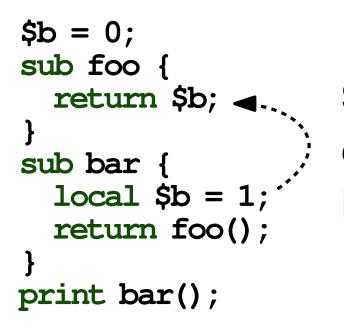
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```
$b = 0;
sub foo {
  return $b; 
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sub bar {
  local $b = 1;
  return foo();
}
print bar();
```

Scope of local \$b dynamically extends into invocation of foo

Quiz: Dynamic Scoping

What does this Perl code print? (Hint: Perl uses dynamic scoping for local variables)



Scope of local \$b dynamically extends into invocation of foo

Answer: 1

Quiz: Static Scoping

What does this Python code print? (Hint: Python uses static scoping)

```
\mathbf{x} = \mathbf{a}
def f():
     def q():
           print(x)
     def h():
           g()
           x = "c"
           print(x)
     \mathbf{x} = \mathbf{b}
     h()
     print(x)
f()
print(x)
```

Quiz: Static Scoping

What does this Python code print? (Hint: Python uses static scoping)

```
x = "a"
def f():
    def q():
         print(x) # (1) x in f : "b"
    def h():
         g()
         \mathbf{x} = \mathbf{C}
         print(x) # (2) x in h : "c"
    \mathbf{x} = \mathbf{b}
    h()
    print(x) # (3) x in f : "b"
f()
                   # (4) x in main: "a"
print(x)
```

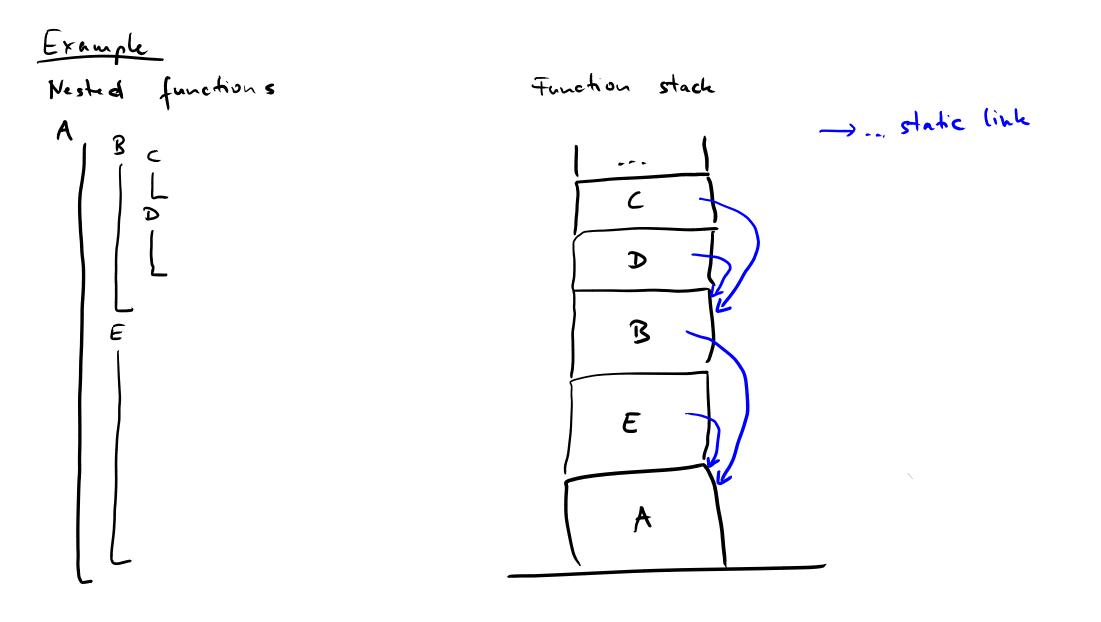
Function Stack vs. Static Scopes

- Push allocation frames on calls
- Pop frames on returns

 Not affected by which functions get called

How to resolve bindings outside of current scope?

 Each allocation frame has a static link to its parent scope



Built-in Objects

Many PLs have built-in (or predefined) objects

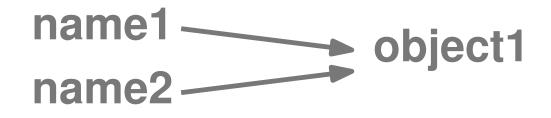
- E.g., for built-in types and APIs
- Invisible, outer-most scope
- Accessible from all scopes, except if hidden

Overview

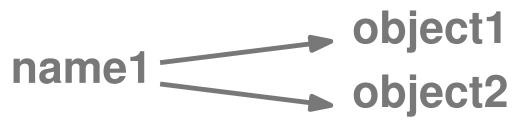
- Object lifetime and storage management
- Scopes
- Aliasing and overloading
- Binding of referencing environments

Aliasing and Overloading

Aliasing: Two more more names refer to the same object



Overloading: A name refers to two more objects



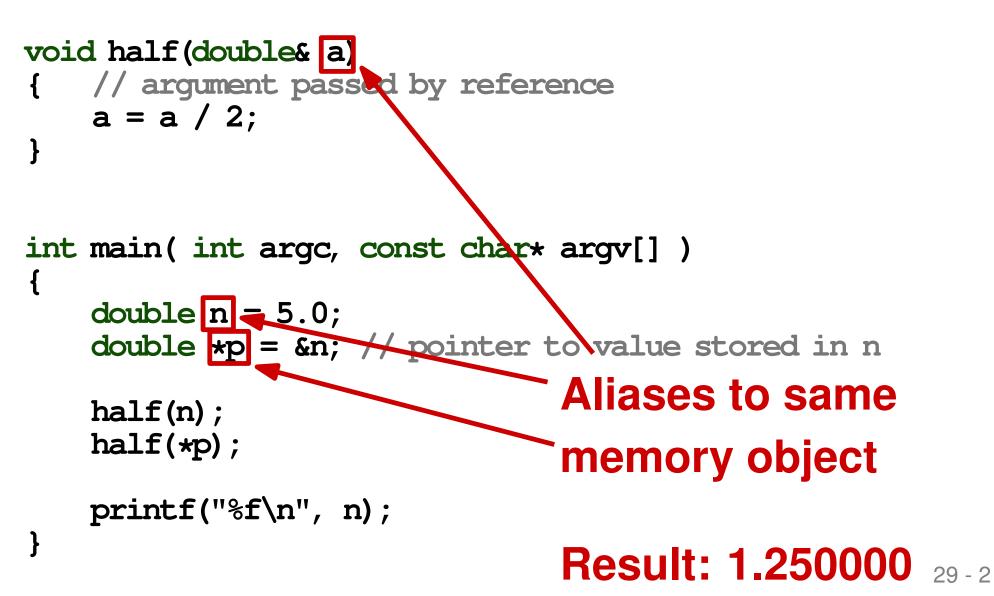
Aliasing: Example

#include <stdio.h>

```
void half(double& a)
{ // argument passed by reference
    a = a / 2;
}
int main( int argc, const char* argv[] )
{
    double n = 5.0;
    double *p = &n; // pointer to value stored in n
    half(n);
    half(*p);
   printf("%f\n", n);
}
```

Aliasing: Example

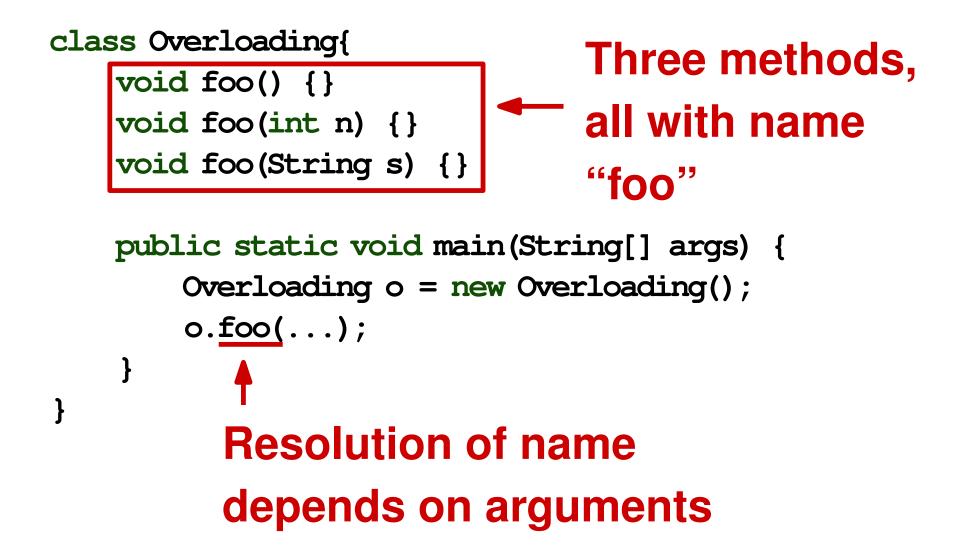
#include <stdio.h>



Overloading: Example

```
class Overloading{
    void foo() {}
    void foo(int n) {}
    void foo(String s) {}
    public static void main(String[] args) {
        Overloading o = new Overloading();
        o.foo(...);
    }
}
```

Overloading: Example



Overview

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- Scopes
- Aliasing and overloading
- Binding of referencing environments

Referencing Environment

Complete set of bindings at a point in the program

 Determined by scoping rules (e.g., static or dynamic scoping)

What if we create a reference to a function?

• When to apply the scoping rules?

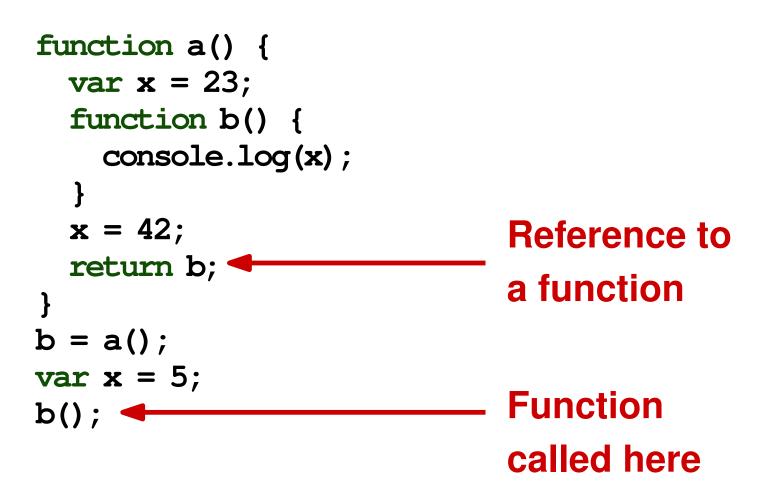
Pseudo code:

```
function a() {
  var x = 23;
  function b() {
    console.log(x);
  }
  x = 42;
  return b;
}
b = a();
var \mathbf{x} = 5;
b();
```

Pseudo code:

```
function a() {
  var x = 23;
  function b() {
    console.log(x);
  }
  x = 42;
                              Reference to
  return b;
                              a function
}
b = a();
var x = 5;
b();
```

Pseudo code:



Pseudo code:

```
function a() {
  var x = 23;
  function b() {
    console.log(x);
  x = 42;
  return b;
}
b = a();
var \mathbf{x} = 5;
b();
```

What memory object is x bound to?

Shallow Binding

Referencing environment created when function is called

Common in languages with dynamic scoping

Shallow Binding

Referencing environment created when function is called

Common in languages with dynamic scoping

```
function a() {
    var x = 23;
    function b() {
        console.log(x);
    }
    x = 42;
    return b;
}
b = a();
var x = 5;
b();
```

Shallow Binding

Referencing environment created when function is called

Common in languages with dynamic scoping

```
function a() {
    var x = 23;
    function b() {
        console.log(x);
    }
    x = 42;
    return b;
}
b = a();
var x = 5;
b();
```

x bound to the global variable initialized to 5; code prints 5

Deep Binding

Referencing environment created when the reference to the function is created

Common in languages with static scoping

Deep Binding

Referencing environment created when the reference to the function is created

Common in languages with static scoping

```
function a() {
    var x = 23;
    function b() {
        console.log(x);
    }
    x = 42;
    return b;
}
b = a();
var x = 5;
b();
```

Deep Binding

Referencing environment created when the reference to the function is created

Common in languages with static scoping

```
function a() {
    var x = 23;
    function b() {
        console.log(x);
    }
    x = 42;
    return b;
}
b = a();
var x = 5;
b();
```

x bound to the local
variable initialized to 23;
code prints 42, as this is
the most recent value of x

Closure

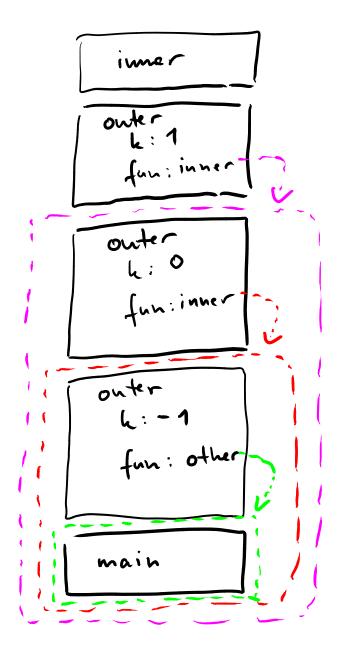
- Implementation of deep binding
- Closure = Representation of referencing environment + function itself
- When creating reference to function, closure is created

Example: Closures

```
function outer(k, fun) {
  function inner() {
    console.log(k);
  }
  if (k > 0)
    fun();
  else
    outer(k + 1, inner)
}
```

function other() {}

```
outer(-1, other);
```



> referencing > environments captured by closures ---> --->

prints 0

Quiz: Scopes and Bindings

Which of the following statements is true?

- Statically allocated memory changes on every function call.
- With nested scopes, a name is resolved using the closest scope that offers a binding for the name.
- Pointers are one way of creating aliases.
- The referencing environment is always created when a function is declared.

Quiz: Scopes and Bindings

Which of the following statements is true?

- Statically allocated memory changes on every function call.
- With nested scopes, a name is resolved using the closest scope that offers a binding for the name.
- Pointers are one way of creating aliases.
- The referencing environment is always created when a function is declared.

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